

# Counting colorings of a regular graph

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# An enumeration question for colorings

## (Proper) $q$ -colouring of graph $G$

- $C : V(G) \rightarrow \{1, \dots, q\}, xy \in EG \Rightarrow C(x) \neq C(y)$
- $c_q(G) =$  number of  $q$ -colourings

## An enumeration question

- Fix  $q$  and family  $\mathcal{G}$  of graphs
- Which  $G \in \mathcal{G}$  maximizes  $c_q(G)$ ?

## $\mathcal{G}$ the family of $n$ -vertex, $m$ -edge graphs

- Question raised by Wilf and Linial
- Partial results, but no complete answer or conjecture

## Our family of interest

$\mathcal{G}(n, d)$  the family of  $n$ -vertex,  $d$ -regular graphs

**Conjecture** (G., Tetali) For  $G \in \mathcal{G}(n, d)$  and all  $q$

$$c_q(G) \leq c_q(K_{d,d})^{\frac{n}{2d}}$$

### Remarks

- Tight when  $2d|n$  ( $G$  the disjoint union of  $n/2d$  copies of  $K_{d,d}$ )
- Conjecture true for *bipartite*  $G \in \mathcal{G}(n, d)$  (G.-Tetali)
- Conjecture true for  $q \geq n^8$  (G., Lazebnik, Zhao)

# An asymptotic approach

An estimate for the conjectured upper bound,  $q = 4$

$$\begin{aligned}c_4(K_{d,d}) &= 6 \cdot 4^d + O(3^d) \\c_4(K_{d,d})^{\frac{n}{2d}} &= 2^n 6^{\frac{n(1+o_d(1))}{2d}}\end{aligned}$$

Best known upper bound to date for  $G \in \mathcal{G}(n, d)$

$$c_4(G) \leq 2^n 6^{\frac{3n(1+o_d(1))}{4d}}$$

Theorem (G.) For  $G \in \mathcal{G}(n, d)$ ,

$$c_4(G) \leq 2^n 6^{\frac{n(1+o_d(1))}{2d}}$$

- Similar bounds for all fixed  $q \geq 3$
- In all cases correct in first two terms of exponent

# Two sides to the coin

## Two regimes in which we work

- $G$  close to bipartite (there's an independent set of size  $\approx n/2$ )
- $G$  far from bipartite (no large independent set)
- Independent set: set of mutually non-adjacent vertices

**Theorem** (G.) If  $G$  has independent set of size  $n(1 - \varepsilon)/2$

$$c_4(G) \leq 2^n 6^{\frac{n(1+\varepsilon)}{2d}}$$

- Uses entropy, and ideas from Kahn, Madiman-Tetali

**Theorem** (G.) If  $G$  has *no* independent set of size  $n(1 - \varepsilon)/2$

$$c_4(G) \leq 2^n \exp \left\{ O \left( n \sqrt{\frac{\log d}{d}} \right) - \Omega(\varepsilon n) \right\}$$

- Take  $\varepsilon = C \sqrt{\log d/d}$  for large enough  $C$  to get the main result

# A tool for dealing with small independent sets

**Lemma** (Sapozhenko) For each  $G \in \mathcal{G}(n, d)$  there is  $\mathcal{F} \subseteq 2^{V(G)}$  with

- Each independent set  $I$  satisfies  $I \subseteq F$  for some  $F \in \mathcal{F}$
- Each  $F \in \mathcal{F}$  satisfies  $|F| \leq \frac{n}{2} \left( 1 + O \left( \sqrt{\frac{\log d}{d}} \right) \right)$
- $\mathcal{F}$  satisfies  $|\mathcal{F}| \leq \exp \left\{ O \left( n \sqrt{\frac{\log d}{d}} \right) \right\}$

## Proof sketch

- Build “seed”  $S$  inside each  $I$ , by repeatedly adding vertices with  $\geq \sqrt{d \log d}$  new neighbours
- Build “cover”  $F$  from  $S$  by taking all vertices not in  $N(S)$  that have many neighbours to  $N(S)$
- Seeds are small, so few seeds and few covers
- Construction of  $F$  ensures it really covers  $I$
- $N(S)$  largish, so  $F$  smallish

# Counting 4-colourings of regular graphs

## A weak bound

- Colouring  $C$  partitions  $V(G)$  into 4 independent sets  $(C_1, C_2, C_3, C_4)$
- $(C_1, C_2, C_3, C_4)$  “covered” by  $(F_1, F_2, F_3, F_4)$
- At most  $\exp \left\{ O \left( n \sqrt{\frac{\log d}{d}} \right) \right\}$  covering 4-tuples
- Number of ways of reconstructing colouring from covering 4-tuples is

$$\prod_{v \in V(G)} a_v \quad (a_v = \text{number of } F_i \text{ in which } v \text{ sits})$$

- AM-GM inequality:

$$\prod_{v \in V(G)} a_v \leq \left( \frac{1}{n} \sum_{v \in V(G)} a_v \right)^n \leq 2^n \exp \left\{ O \left( n \sqrt{\frac{\log d}{d}} \right) \right\}$$

- Conclusion: for  $G \in \mathcal{G}(n, d)$ ,  $c_4(G) \leq 2^n \exp \left\{ O \left( n \sqrt{\frac{\log d}{d}} \right) \right\}$

# Incorporating the independent set information

Knowing  $G$  has no independent set of size  $n/2(1 - \varepsilon)$

- May assume all  $F_i$ 's have at least  $n/2$  vertices
- $F_1$  has no large independent sets so has  $\Omega(\varepsilon n)$ -sized matching
- For  $uv$  in matching,  $a_u a_v$  in reconstruction replaced by  $a_u a_v - 1$
- Dampens the count sufficiently to give main result

## Final remarks

- All ok for even  $q$ ; for odd  $q$  need to use that  $a_v$ 's are integers
- Proof of tight bound remains to be found
- We've shown: for all  $\varepsilon > 0$  there's  $\delta > 0$  so that for  $G \in \mathcal{G}(n, d)$  without independent sets of size  $n(1 - \varepsilon)/2$ ,

$$c_4(G) \leq (2 - \delta)^n$$

How does  $\delta$  depend on  $\varepsilon$ ?