

Graphic Representations  
of  
Long Chordless Cycles

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**host graph**  $H$  where the representation occurs

**target graph**  $G$  that we wish to represent

We assign a representative subgraph  $R_v$  of  $H$  to each vertex  $v$  in  $G$  so that  $v$  and  $w$  are adjacent iff  $R_v$  and  $R_w$  “conflict” — that is, they have “enough in common”.

Examples:

Host: A tree  $T$

Targets: Chordal graphs

Representatives: Subtrees of  $T$

“Enough” (Conflict rule): A node

NOTE: Vertices of the target are “vertices”  
Vertices of the host are “nodes”.

In this talk, all representatives will be isomorphic to a fixed **representative prototype**.

The conflict rule kicks in when two representatives contain a common copy of a fixed **quota**.

Host: A graph  $H$

Targets: Line graphs

Representatives:  $P_2$  — i.e., edges of  $H$

Quota:  $P_1$  — i.e., a node of  $H$ .

Representative prototype:  $P_4$

Quota prototype:  $P_2$

$r$ : order of the representative prototype  $R$   
 $q$ : order of the quota prototype  $Q$

An  $(H; R, Q)$ -*representation* of a graph  $G$  in a host  $H$  is an injective assignment  $v \rightarrow R_v$  of a copy  $R_v$  of  $R$  to each vertex  $v$  of the target such that

$$\begin{aligned}vw \in E(G) \\ \iff \\ R_v \cap R_w \text{ contains a copy of } Q.\end{aligned}$$

The *universal graph* has all copies of  $R$  in  $H$  as vertices, with adjacency determined as above.

$G$  is representable  $\iff$

$G$  an induced subgraph of the universal graph

Let  $K_n$  be the host.

Let  $M$  be the maximum number of  $Q$ -copies contained in an induced subgraph of  $H$  of order  $q$ .

Let  $G$  be the target graph with maximum clique size  $\omega$ .

The order of  $G$  is bounded by  $\omega M \binom{n}{q}$ .

$M$  and  $\omega$  are constants with respect to  $n$ , so the maximum order of representable graphs with bounded clique size is  $O(n^q)$ .

This order is achieved if the representative prototype is  $P_r$  and the quota is  $P_q$ .

Assumptions:

$$r > q \geq 2$$

For a fixed  $m \geq 2$  assume

$r - q$  and  $r$  both divide  $m^q$ .

Select a de Bruijn sequence  $a_k$  on an alphabet of  $m$  letters

such that each  $q$ -tuple occurs exactly once.

Take the alphabet to be  $[m]$ .

Reference: Robert E. Jamison, Towards a Comprehensive Theory of Conflict-Tolerance Graphs, Proceedings of LAGOS Conference, to appear in Discrete Applied Math.