

# Finding the Kth largest item in a list of n items

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# Content

- 1 Introduction
- 2 History
- 3 Comparison
- 4 Algorithm
- 5 Example
- 6 Performance
- 7 Applications

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- This problem can be solved by two algorithms:
  - Selection algorithm.
  - Median of medians algorithm.
- Median of medians algorithm is better than selection algorithm due to its worst case linear time performance.

# History

The Median of Medians Algorithm was proposed by 5 great computer scientists they are **Manuel Blum, Robert Floyd, Vaughan Pratt, Ron Rivest** and **Robert Tarjan** in the year **1973**.



# Comparison

- In median of medians algorithm, we divide the list by 5 and then we sort the divided list, where as in selection algorithm we directly sort the unordered list with out dividing.

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- In median of medians algorithm, we divide the list by 5 and then we sort the divided list, where as in selection algorithm we directly sort the unordered list with out dividing.
- Median of medians algorithm has a better performance when compared to selection algorithm.

# Algorithm

- 1 Divide the list in to  $n/5$  lists of 5 elements each.
- 2 Find the median in each sublist of 5 elements.
- 3 Recursively find the median of all the medians, call it  $m$ .
- 4 Partition the list in to unique elements larger than ' $m$ ' (call this sublist  $L1$ ) and those no longer than ' $m$ ' (call this sublists  $L2$ ).
- 5 If  $K \leq |L1|$ , return selection ( $L1, K$ ).
- 6 If  $K - 1 = |L1|$ , return ' $m$ '.
- 7 If  $K > |L1| + 1$ , return selection( $L2, K - |L1| - 1$ ).

# Example

Find the 8th largest element i.e  $k=8$ .

Unordered list

2 3 5 4 1 12 11 13  
16 7 8 6 10 9 17 15  
19 20 18 23 21 22  
25 24 14

Dividing by 5

2 3 5 4 1  
12 11 12 16 7  
8 6 10 9 17  
15 19 20 18 23  
21 22 25 24 14

# Example

## Finding medians

2	3	5	4	1
12	11	13	16	7
8	6	10	9	17
15	19	20	18	23
21	22	25	24	14

## Median of medians

5 13 10 20 25

# Example

**2 3 5 4 1 7 8 6 9**

$|L2|=9$

**10**

**m**

**12 11 13 16 17 15 19 20**  
**18 23 21 22 25 24 14**

$|L1|=15$

# Example

12 11 **13** 16 17  
15 19 **20** 18 23  
21 22 **25** 24 14

→ 13 **20** 25

12 11 13 16 17 15 19  
18 14

$|L2|=9$

**20**

23 21 22 25 24

$|L1|=5$

# Example

12 11 13 16 17  
15 19 18 14

12 11

13

14 15 16 17 18 19

$|L1| = 6; K = 2$



# Example

14 15

16

17 18 19

$L1=3; k=2$

17

18

m

19

$|L1|=1; k=2$

## Why 5?

- Dividing the list by 5 assures a worst-case split of 70 – 30.
- At least half of the medians are greater than the median-of-medians, hence at least half of the  $n/5$  blocks have at least 3 elements and this gives a  $3n/10$  split, which means the other partition is  $7n/10$  in worst case.

That gives  $T(n) = T(n/5) + T(7n/10) + O(n)$ .

# Performance

- The best total running time of finding *Kth* largest item in a list of *N* items is  $O(n \log n)$ .
- Where running time of sorting *N* items is  $O(n \log n)$  and running time of returning the *Kth* largest item is  $O(1)$ .
- The worst case running time is  $O(n)$ .

# Applications

- **Order Statistics:** Selection include for finding the smallest elements, largest elements and median elements.
- **Computer chess program:** Identifying the most promising candidates in computer chess program.
- **Salary Distribution:** Selection is used in salary distribution.
- **Filtering Outlying** elements such as noisy data.

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